

# Modeling Task Allocation Using a Decision Theoretic Model\*

Sherief Abdallah and Victor Lesser  
University of Massachusetts, Amherst  
MAS Laboratory

{shario,lesser}@cs.umass.edu

## ABSTRACT

Mediation is the process of decomposing a task into subtasks, finding agents suitable for these subtasks and negotiating with agents to obtain commitments to execute these subtasks. This process involves several decisions to be made by a mediator including which tasks to mediate, when to interrupt the current task mediation to pursue a better task, etc. The main contribution of this work is integrating the different aspects of a mediator decision problem into one coherent and simple decision theoretic model. This model is then used to learn an optimal policy for a mediator.

We propose a generalization of the original Semi-MDP (SMDP) model, which allows efficient representation of the mediator decision problem. Also the concurrent action model (CAM) is extended to allow better performing policies to be found. Experimental results are presented showing how our model outperforms the original SMDP and CAM models.

## Categories and Subject Descriptors

I.2.6 [Artificial Intelligence]: Learning; I.2.11 [Artificial Intelligence]: Distributed Artificial Intelligence

## General Terms

Algorithms, Experimentation

## Keywords

Reinforcement Learning, Multiagent Systems, Task Allocation, Markov Decision Process

## 1. INTRODUCTION

Task allocation is a problem common to most multi-agent systems. An agent discovers a task that it can not accomplish by itself. The

\*This material is based upon work supported [in part] by the National Science Foundation Engineering Research Centers Program under NSF Award No. EEC-0313747. Any opinions, findings, conclusions or recommendations expressed in this material are those of the author(s) and do not necessarily reflect the views of the National Science Foundation.

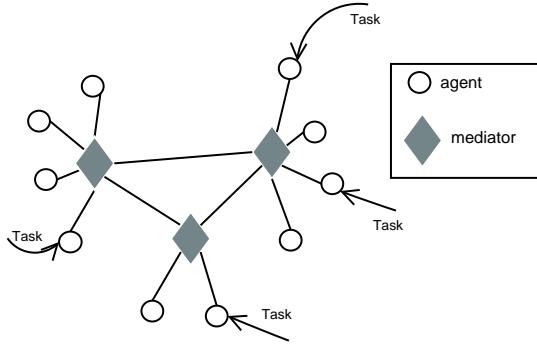
agent then finds other agents with the right capabilities and assigns the task, or parts of it, to those agents. We use the term *mediation* to refer to the process of decomposing a task into subtasks, finding agents suitable for these subtasks and negotiating with agents to obtain commitments to execute these subtasks. This process involves several decisions to be made by a mediator including which tasks to mediate (when more than one task is available), when to interrupt the current task mediation to pursue a better task, which agent to assign to which subtask, etc. Mediation does not, necessarily, sacrifice an agent's autonomy for efficiency. A mediator may assign a task to an agent but this agent may reject this assignment if there is a more profitable course of action.

The main distinguishing characteristic of a mediator decision problem is that tasks appear randomly. This means the set of actions available to a mediator (associated with these tasks) changes *non-deterministically* from time to time. We refer to this property as the *randomly available actions* property. As we show in this paper, existing MDP models fail to exploit this property, which makes existing models both non-scalable and suboptimal. The main contribution of this work is integrating the different aspects of a mediator decision problem into one coherent and simple decision theoretic model. In doing so we extend existing SMDP models to exploit the *randomly available actions* property. We then show how our model can be used to learn a mediator's policy that outperforms policies learned using existing models.

The details of the underlying negotiation protocol are not the focus of this work. A mediation process will be abstracted by an *action* with non-deterministic behavior. All that matters is the outcome of that action. For example, the contract net protocol [6] can be used as an underlying negotiation protocol. A mediator's action in that case would be to contract a task to an agent. What matters, from this work's perspective, are the final outcomes of that action: its duration, the costs incurred, and whether it succeeded or failed. This work does not worry about the fine details of the contract net protocol (e.g., sending a bid, receiving a counter bid, and so on until the negotiation either succeeds or fails). Instead, the focus is on the higher-level decision of whether to start contracting a specific task and when to abort mediating a high-level task even if some subtasks have already been committed. Note that from our model's perspective, it does not matter whether an action abstracts just the mediation of a task or the completion of both the task's mediation and then the task itself.

To more deeply understand the decisions that a mediator needs to make, let us first revisit the mediation process in more detail. A mediator receives task announcements. Each announced task has

an associated payoff. Tasks can be decomposed into a set of subtasks. Subtasks are independent from one another but all subtasks must succeed for a mediator to receive the promised payoff. Different agents have different capabilities, hence subtasks may take different durations (depending on the agents assigned to them). The duration is also stochastic (though its mean may depend on the assigned agent). Subtasks also incur costs that need to be covered by the overall payoff for the total task to be *profitable*. The goal of the mediator is to maximize its net profit (maximizing payoff and minimizing cost). In general, mediators may interact with each other in the system, as shown in Figure 1. However, from a mediator's perspective, it receives task announcements and contracts out subtasks to neighboring agents, independent of whether the other agents are mediators or not.



**Figure 1:** A network of mediators for assigning agents to tasks.

Although this work's focus is on formalizing the decision process of a single mediator, this is done in a multi-mediator context. In particular, when a mediator  $m_1$  asks another mediator  $m_2$  to do a task  $T_1$ ,  $m_2$  needs to decide as soon as possible whether to accept or reject  $T_1$  so that  $m_1$  can also decide its course of action. Consequently, a mediator cannot defer its decision about a task and queue it for consideration later. This constraint makes the mediator decision problem more challenging as the mediator needs to make a decision now that may affect its future profit. With this in mind, our work provides a formal decision theoretic model that enables a mediator to evaluate quantitatively the following three decisions.

- How a mediator should select among available tasks, where their subtasks are contracted out *concurrently*. The mediator should take into account that subtasks may have different durations depending on the agents assigned to the subtasks. Furthermore, no reward will be given until all subtasks are complete.
- The mediator needs to decide between *committing now* to a task or *waiting* for a better task to arrive in the future. This decision is particularly important if the mediator cannot de-commit a task once committed. The mediator needs to take into account the payoff of currently available tasks, their duration, and the likelihood of a better task arriving *before* any of the current tasks finish.
- If decommitment is allowed, and the mediator receives a new task announcement, it needs to decide whether it is better to continue with an old task or decommit from the old task and start a new one. This decision needs to take into account the costs already accrued for performing the old task, the payoff for its completion, the expected time for the old task to finish,

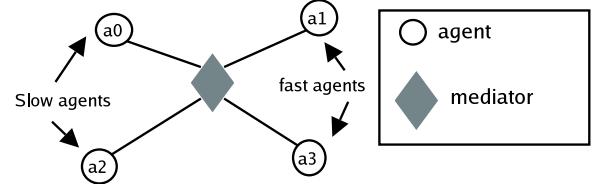
the payoff for the newer task and its expected duration and cost.

The decisions above are further complicated by the fact that subtasks' costs and durations are stochastic with a distribution that is not known to the mediator *a priori*. We show how to model these decisions efficiently by generalizing the Semi-MDP model (SMDP)[8] and extending the Concurrent Action Model (CAM)[5]. We also present an algorithm that uses this model and reinforcement learning techniques to learn an optimal policy for a mediator.

Section 2 presents an example to illustrate the mediator decision problem. Section 3 briefly describes the SMDP model and how it can be used to model serialized mediation. Section 4 shows how the SMDP model can be generalized to efficiently model randomly available actions. Section 5 describes how CAM can be extended to model concurrent mediation of subtasks. Section 6 shows experimental results illustrating how our model outperforms existing models when used to learn an optimal policy.

## 2. EXAMPLE

Let us consider a simple mediator example illustrated in Figure 2. A mediator  $m$  can negotiate with four agents  $a_0, a_1, a_2$ , and  $a_3$ . Each agent  $a_i$  can execute a single subtask  $T_{a_i}$ . The duration and the cost of a subtask  $T_{a_i}$  are stochastic and unknown to the mediator (but they depend on agent  $a_i$ ). They are represented by random variables  $D_{a_i}$  and  $C_{a_i}$  respectively. At any given time, there is a set of available tasks  $T = \{T_0, \dots, T_n\}$  available to  $m$  to choose from to mediate. Each task  $T_k$  can be decomposed into a set of subtasks  $T_k = \{T_{a_i}, T_{a_j}, \dots\}$  and has a promised reward  $R_k$  (both are known to the mediator).



**Figure 2:** The experiment scenario

The probability that a task of type  $T_k$  is available for mediation (at any point of time) is  $p_{T_k}$ . If a task type is not selected for mediation it may or may not be available at the next time step depending on  $p_{T_k}$ . For simplicity, assume the mediator can only mediate one task at a time. Consider an example scenario involving two tasks  $T_0 = \{T_{a_0}, T_{a_1}\}$  and  $T_4 = \{T_{a_1}, T_{a_3}\}$ . At time 0 task  $T_0$  is the only available task.  $m$  starts mediating  $T_0$  with agents  $a_0$  and  $a_1$ . At time 5 mediation about  $T_{a_1}$  finishes with cost  $c_{a_1}$ . At time 7 another task  $T_4$  arrives while the mediation about subtask  $T_{a_0}$  has not finished yet. What should  $m$  do?

If decommitment is allowed then  $m$  needs to consider whether to drop  $T_0$  and start  $T_4$ . Ignoring the effect of duration,  $m$  should drop  $T_0$  if the profit of  $T_4$  is higher, i.e.,  $R_0 - c_{a_1} - E\{C_{a_0}\} < R_4 - E\{C_{a_1}\} - E\{C_{a_3}\}$ , where  $E\{X\}$  is the expectation of random variable  $X$ . The above decision rule is not actually optimal as it ignores the effect of duration. For example, if the expected duration of  $T_{a_3}$  is 20 while the expected duration of  $T_{a_0}$  is 10, it would be better for  $m$  to stick with  $T_0$  even if the profit is lower as this will free  $m$  to accept another task earlier.

Now consider if  $m$  is not allowed to decommit a task once it started mediation. The mediator still has an interesting decision to make: whether to wait for a better task or not. For example in the previous scenario, if  $E\{D_{a_0}\} = 30$  while  $E\{D_{a_3}\} = 5$ , then  $m$  would be better off not accepting  $T_0$  even if at the time  $m$  has nothing else to do and wait for  $T_4$  to arrive. Clearly making such a decision depends on the arrival rate of different tasks. For example, if  $T_4$  tasks arrives on average every 100 time units,  $m$  may be better not waiting for it, while if it arrives every 10 time units it might be worth the wait.

Integrating all these factors into one coherent and simple model is a challenging task. The next section describes the Semi-MDP model which we use as a basis for modeling the mediator's decision.

### 3. SEMI MARKOV DECISION PROCESS, SMDP

The Semi-MDP [8], or SMDP, provides a basis for modeling an agent's decision if the agent's actions may take different durations. An SMDP is defined by the tuple  $\langle S, A, P, R \rangle$ .  $S$  is the set of states,  $A(s)$  is the set of actions available at a given state,  $P(s, a, N, s')$  is the probability of reaching state  $s'$  after  $N$  time-steps of executing action  $a$  at state  $s$ , and  $R(s, a)$  is the average reward if action  $a$  is executed at state  $s$ .

The solution to an SMDP decision problem is an optimal policy  $\pi^*(s)$  that specifies for every state  $s$  what is the best action to execute. The optimal value of action  $a$  at state  $s$ ,  $Q^*(s, a)$  is the expected sum of rewards if the agent first execute action  $a$  then follows the optimal policy thereafter.  $Q^*(s, a)$  is given by Bellman's equation [8]:

$$Q^*(s, a) = R(s, a) + \sum_{N, s'} P(N, s'|s, a) \gamma^N \max_{a'} Q^*(s', a')$$

where  $\gamma$  is a learning parameter called the discount factor. Another variant of this equation which is used in the Q-learning algorithm [7] is  $Q^*(s, a) \leftarrow (1 - \alpha)Q^*(s, a) + \alpha[R(s, a) + \gamma^N \max_{a'} Q^*(s', a')]$ , where  $\alpha$  is another learning parameter called the learning rate. Once the values of  $Q^*$  are learned, the optimal policy can be easily computed:  $\pi^*(s) = \text{argmax}_a Q^*(s, a)$ .

The mediator decision problem can be formulated using an SMDP as follows. The state  $s$  would be a combination of the states of its neighboring agents (or some abstraction of it) and the status of task(s) the mediator is pursuing. An action corresponds to the process of assigning a subtask to an agent. The task payoff is the reward received when the status of the pursued task is *completed successfully*. This is true iff all subtasks complete successfully. Every time-step a subtask may produce a reward  $r$  that is usually (but not necessarily) negative to represent the cost of doing the subtask.

The SMDP model above can handle serialized contracting of sub-tasks. It can correctly model mediations with variant durations. However, this model does not explicitly handle concurrent actions and therefore can not handle concurrent mediation of subtasks. Mediating subtasks in parallel means some of them may finish earlier than others (since subtasks have variant durations). Shall the mediator wait for other subtasks to finish before making a new decision or shall it reconsider its decision whenever a subtask finishes? CAM [5], addresses these issues for the Semi-MDP model with concurrent actions. Section 5 briefly describes CAM and how to extend it to account for stochastically changing set of avail-

able actions. The following section proposes a generalization of the SMDP model that enables more efficient representation of the class of decision problems where the set of actions changes non-deterministically (i.e., the randomly available actions property).

### 4. RANDOMLY AVAILABLE ACTIONS

In the SMDP model, the set of available actions  $A_t$  at time  $t$  is defined by a *deterministic* function  $A(s_t)$  where  $s_t$  is the state at time  $t$ . In other words, the SMDP model requires the set of available actions at any time to be *deterministically* defined by the agent's state.

In the task allocation decision problem, the set of available actions  $A_t$  corresponds to the set of available tasks at time  $t$ . Because tasks arrive in a non-deterministic fashion, the only way to satisfy the SMDP model requirement of having a deterministic function  $A(s)$  is to include the set of available tasks as part of the state. While compatible with the traditional SMDP, this results in an exponential explosion in the number of states that hinders learning.

However, it is intuitive that in some domains the value of an action is independent of other available actions (we prove this below). However, requiring the set of available actions to be a deterministic function of the current state is over constraining. We propose a generalization of the SMDP model,  $\wp$ -SMDP, which allows the set of available actions  $A_t$  at time  $t$  to be *non-deterministically* dependent on the current state of the agent. As we show below, this generalized model allows exponential savings in the state space.

$\wp$ -SMDP is defined by the tuple  $\langle S, \wp, P, R \rangle$ , where  $S$ ,  $P$ , and  $R$  are as described in the previous section, while  $\wp(A|s)$  is the probability that the set of actions  $A$  is available at state  $s$ . The  $\wp(A|s)$  replaces the traditional  $A(s)$  function allowing more flexibility in the model. Mapping a traditional SMDP to  $\wp$ -SMDP is straightforward.  $S$ ,  $P$ , and  $R$  are the same.  $\wp(A(s)|s) = 1$  and  $\forall A \neq A(s) \wp(A|s) = 0$ . For clarity, we will use the symbol  $\wp$  to refer to probabilities associated with actions. For example,  $\wp(A^{t+N} = A' | s^{t+N} = s', s^t = s)$  is the probability that the set of available actions at time  $t + N$  is  $A$  given that the state at time  $t + N$  is  $s'$  and the state at time  $t$  is  $s$ . We will also omit random variables when possible, i.e.  $\wp(A^{t+N} = A' | s^{t+N} = s', s^t = s) = \wp(A' | s', s)$ .

The real benefit of  $\wp$ -SMDP over the original SMDP is not its generality, but that under certain assumptions  $\wp$ -SMDP leads to a more compact representation of the Bellman equations as we show below. In particular, it is possible in the  $\wp$ -SMDP to factor the set of available actions out of the state. This leads to *exponential* savings in the state space and hence reduces the memory and time required to learn an optimal policy.

Let  $s$  be the state with the set of available actions factored out while  $\langle A, s \rangle$  be the state with the set of available actions factored in. Let  $A$  and  $A'$  be the set of available actions at time  $t$  and  $t + N$  respectively. Also  $s$  and  $s'$  are the state at  $t$  and  $t + N$  respectively; and  $a$  and  $a'$  are actions at  $t$  and  $t + N$ .

**THEOREM 1.** : If the following is true:

- *action independence:* the set of available actions at time  $t$  can be dependent only on the current state, and is independent of all other variables given the current state. So for example,  $\wp(A'|N, s', \langle A, s \rangle, a_i) = \wp(A'|s')$ .

- *outcome independence:* the outcome of applying an action  $a$  in a given state  $s$  does not depend on the set of available actions at that time, i.e.  $P(N, s'|A, s, a) = P(N, s'|s, a)$

then the equation below is also true:

$$Q(\langle A, s \rangle, a_i) = Q(s, a_i) = R(s) + \sum_{N, A', s'} \wp(A'|s') P(N, s'|s, a_i) \gamma^N \max_{a'_i \in A'} Q(s', a'_i)$$

PROOF. Using Bellman's equation,

$$Q(\langle A, s \rangle, a_i) = R(s) + \sum_{N, A', s'} P(N, \langle A', s' \rangle | \langle A, s \rangle, a_i) \gamma^N \max_{a'_i \in A'} Q(\langle A', s' \rangle, a'_i)$$

but, using the Probability Product rule,

$$\begin{aligned} P(N, \langle A', s' \rangle | \langle A, s \rangle, a_i) &= \\ \wp(A'|N, s', \langle A, s \rangle, a_i) P(N, s' | \langle A, s \rangle, a_i) \wp(A'|N, s', \langle A, s \rangle, a_i) \\ &\quad P(N, s' | \langle A, s \rangle, a_i) \end{aligned}$$

then, using the action independence assumption for the first part and the outcome independence for the second part,

$$= \wp(A'|s') P(N, s' | s, a_i)$$

and hence

$$Q(\langle A, s \rangle, a_i) = R(s) + \sum_{N, A', s'} \wp(A'|s') P(N, s' | s, a_i) \gamma^N \max_{a'_i \in A'} Q(\langle A', s' \rangle, a'_i) \quad (1)$$

where parameter  $A$  in  $Q(\langle A, s \rangle, a_i)$  does not appear in the r.h.s., therefore  $Q(\langle A, s \rangle, a_i) = Q(s, a_i)$ . Thus we can substitute  $Q(\langle A', s' \rangle, a'_i) = Q(s', a'_i)$  in (1).  $\square$

Note that the action-independence and outcome-independence assumptions are not as restrictive as it may first appear. The reason is that we did not impose any restriction on the state definition, allowing trade-off between compactness and accuracy. For example, in many domains the set of available actions may include the action the agent chose previously and is still executing (because the agent may choose either to continue the currently executing action or start a new action). In this case, it is sufficient to include the currently executing action as part of the state.

The theorem above leads to compact representation of action values by factoring out the set of available actions. In other words, the values of all states that only differ in the set of available tasks are considered equal. This may lead, as the experiments show, to faster convergence using either dynamic programming methods (e.g. value iteration) or temporal difference methods (e.g. Q-learning). For completeness, Algorithm 1 shows how the value iteration algorithm [7] can be adopted (however, this work focuses on Q-learning).

The corresponding update equation for Q-learning is  $Q(s, a) = (1 - \alpha)Q(s, a) + \alpha(R(s, a) + \gamma^N * \max_{a' \in A'} Q(s', a'))$  and the optimal policy  $\pi^*(A, s) = \operatorname{argmax}_{a \in A} Q(s, a)$ . Because the set of available actions is factored out of the state, the optimal action at any state is not unique. The optimal action depends also on the set

---

#### Algorithm 1: Value Iteration

---

```

begin
  Initialize  $Q(s, a)$  arbitrarily,  $\forall a, s$ .
  repeat
     $\delta \leftarrow 0$ 
    for every state  $s$  and action  $a$  do
       $q \leftarrow Q(s, a)$ 
       $Q(s, a) \leftarrow R(s, a) +$ 
       $\sum_{N, A', s'} \wp(A'|s') P(N, s' | s, a) \gamma^N \max_{a' \in A'} Q(s', a')$ 
       $\delta \leftarrow \max(\delta, |Q(s, a) - q|)$ 
    end
  until  $\delta > \epsilon$  (small positive number)
end

```

---

of available actions at that time. However, the optimal policy need not be stored explicitly. Instead, the optimal action at time  $t$  can be dynamically computed using  $Q(s^t, .)$  and  $A^t$ , where  $s^t$  and  $A^t$  is the state and set of available actions at time  $t$  respectively.

Theorem 1 does not violate the conditions for convergence for Q-learning. In particular, Q-learning will converge with enough sampling of states and set of available actions, appropriately decaying learning rate, and bounded rewards [3]. Also since SMDP is a more generic model than MDP, the theorem holds for MDP.

#### 4.1 The *wait* operator

Having a dynamic set of available actions that stochastically changes over time raises an interesting decision question: whether an agent (mediator) should *wait* for a better action (task) to be available in the future or should the agent choose the best action (task) that is available now. One would expect that the probability distribution of the availability of actions (tasks),  $\wp$ , should affect such a decision. Our experiments in Section 6 verify this intuition.

### 5. EXTENSION TO CONCURRENT ACTION MODEL

The Concurrent Action Model (CAM) [5] allows concurrent actions with different durations. The model builds on the Semi-MDP model. The action value equation of CAM is almost identical to that of the original SMDP model:

$$Q(s, \bar{a}) = R(s, \bar{a}) + \sum_{N, s'} P(s, \bar{a}, N, s') \gamma^N \max_{\bar{a}'} Q(s', \bar{a}')$$

where the only difference is that  $\bar{a}$  here is a set of concurrent actions (also called *joint actions*) instead of a single action. Thus using this model a mediator can contract out all subtasks in parallel, and this is considered a single joint action.

The main contribution of [5] is the introduction of *termination schemes* which determine when the next decision cycle occurs. They defined three *termination schemes*:  $\tau_{any}$ ,  $\tau_{all}$ , and  $\tau_{continue}$ . The termination scheme  $\tau_{any}$  means that an agent makes a new decision if any concurrent action finishes, possibly decommitting some of the unfinished actions (tasks). The termination scheme  $\tau_{continue}$  is similar to  $\tau_{any}$  but without decommitting. Finally, in  $\tau_{all}$  the agent waits until all concurrent actions that started together to finish before making a new decision. It was proved that optimal policies learned using  $\tau_{any}$  dominates both  $\tau_{all}$  and  $\tau_{continue}$  [5].

In our case, the set of available actions may change at any time even if none of the currently executing actions have terminated. None of the termination schemes proposed in [5] account for that.

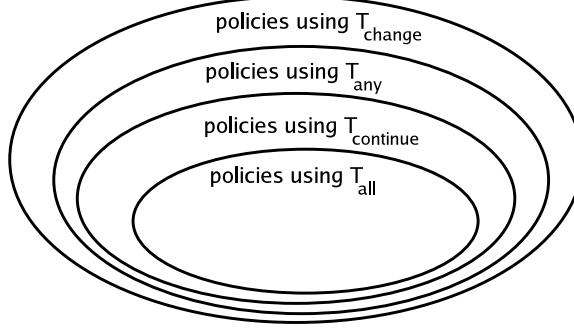
We propose the  $\tau_{change}$  termination scheme, in which a decision starts whenever an action terminates or the current state changes or the set of available actions change. In the task allocation problem  $\tau_{change}$  means a new decision is made when a mediation terminates or a new task arrives.

$\tau_{change}$  is more suitable for multiagent settings where the state observed by an agent is not fully controlled by the actions of that agent. Other agents may change the system's state unexpectedly. Next theorem compares the policy learned under  $\tau_{change}$  to a policy learned under other termination schemes proposed in [5]. Let  $Q^{*\tau}$  be the optimal action values learned using termination scheme  $\tau$ .

**THEOREM 2.** *The optimal policy obtained using  $\tau_{change}$  has a value that is higher than or equal to the value of the optimal policy using  $\tau_{any}$ , or more formally:  $\forall s \in S, \max_{\bar{a}} Q^{*\text{any}}(s, \bar{a}) \leq \max_{\bar{a}} Q^{*\text{change}}(s, \bar{a})$*

**PROOF.** For every policy  $\pi^{*\text{any}}$  obtained using  $\tau_{any}$  termination scheme there exists a corresponding policy  $\pi^{*\text{change}}$  using  $\tau_{change}$  with the same value.  $\pi^{*\text{change}}$  can be obtained from  $\pi^{*\text{any}}$  simply by not changing the executed action unless an action terminates (i.e. repeating the previous action if the decision epoch was initiated due to a change in state or set of available actions).  $\square$

In other words, the space of policies that can be learned using  $\tau_{change}$  dominates that of  $\tau_{any}$  which in turn dominates  $\tau_{continue}$  and  $\tau_{all}$ . Figure 3 augments the figure in [5] with  $\tau_{change}$  policy space. Results shown in Section 6 show that in practice, the policy learned using  $\tau_{change}$  outperforms policies that use other termination schemes.



**Figure 3:** The relationship between policies learned using  $\tau_{all}, \tau_{continue}, \tau_{any}$ , and  $\tau_{change}$

The following section outlines an algorithm for the decision of the mediator. Part of the algorithm learns the optimal policy using the update equation described in Section 4. This algorithm is used in the experiments of Section 6.

### 5.1 Learning the Mediator's Decision Process

Algorithm 2 illustrates how mediators make decisions and learn the values of their actions. The main variables are:  $s$ , the current state,  $s_{last}$ , the previous state, and  $a_{last}^*$ , the action executed at  $s_{last}$ . Steps 2.1-2.4 correspond to Q-learning (using a factored state). Steps 2.5-2.6 are the communication part, which can be modified if more complex protocol is desired. Steps 2.7-2.10 update the main

variables of the algorithm. Algorithm 3 generate actions for the set of available tasks.

---

#### Algorithm 2: Learning

---

```

2.1 begin
2.2   Initialization.  $s \leftarrow$  current state,  $\bar{a}_{last}^* \leftarrow$  null, and
       $s_{last} \leftarrow$  null
2.3   Let  $A \leftarrow \text{generateActions}()$  be the set of actions
      available at this time step
2.4   with probability  $\epsilon$  (small positive number) let
       $\bar{a}^* \leftarrow$  random action  $\in A$  and with probability
       $(1 - \epsilon)$  let  $\bar{a}^* \leftarrow \text{argmax}_{\bar{a} \in A} Q(s, \bar{a})$ .1 Let  $T^*$  be
      the task associated with  $\bar{a}^*$ , and  $R^*$  be the reward
      that will be gained if  $T^*$  is accomplished
2.5   If  $\bar{a}_{last}^* \neq$  null then learn:  $Q(s_{last}, \bar{a}_{last}^*) =$ 
       $(1 - \alpha)Q(s_{last}, \bar{a}_{last}^*) + \alpha[R + \gamma^N Q(s, \bar{a}^*)]$ 
2.6   Send negative response to rejected tasks: send
      reward  $R_{fail}$  to all but the originator of  $T^*$ 
2.7   Start mediating subtasks of  $T^*$  with their assigned
      agents/neighbors as specified by the selected set of
      concurrent actions/assignments  $\bar{a}^*$ 
2.8   Let  $R \leftarrow 0$  and wait for the termination scheme
      condition to be satisfied after  $N$  time-steps (i.e.,
      wait for all subtasks to finish in case of  $\tau_{all}$ , any
      subtask to finish in case of  $\tau_{any}$ , or any subtask to
      finish or any new task to arrive in case of  $\tau_{change}$ )
2.9   For every received response with reward  $r_k$  from a
      neighbor  $k$ ,  $R \leftarrow R + \gamma^N r_k$ 
2.10  If task  $T^*$  is accomplished (i.e. all subtasks
      succeed), then  $R \leftarrow R + R^*$ 
2.11   $s_{last} \leftarrow s, \bar{a}_{last}^* \leftarrow \bar{a}^*$ 
2.12  Goto 2
2.13 end

```

---

### 5.2 Handling Multiple Tasks in Parallel

The concurrent action model can also be used to model a mediator who can mediate multiple tasks (and their corresponding subtasks) in parallel. Another level of joint actions is needed. Sub-actions in that level are tasks which in turn are joint actions of subtask mediations. For example, consider the scenario described in Section 2. Corresponding to tasks  $T_0$  and  $T_4$  the mediator has the joint actions  $\bar{a}_{T_0} = \{a_{T_{a_0}}, a_{T_{a_1}}\}$  and  $\bar{a}_{T_4} = \{a_{T_{a_1}}, a_{T_{a_3}}\}$  respectively, where  $a_{T_{a_i}}$  is the assignment  $\langle T_{a_i}, a_i \rangle$ . To allow the mediator to negotiate about tasks  $T_0$  and  $T_4$  concurrently, another action needs to be added to the pool of available actions:  $\bar{a}_{T_0, T_4} = \{\bar{a}_{T_0}, \bar{a}_{T_4}\}$ .

Figure 4 illustrates the hierarchy of action  $\bar{a}_{T_0, T_4}$ . A different termination scheme can be assigned to each level of this hierarchy of joint actions. For example, the termination scheme for the subtasks level may be  $\tau_{all}$  (wait for all subtasks to finish) while the termination scheme for the tasks level may be  $\tau_{any}$  (if any task finishes, start mediating another one).

## 6. RESULTS

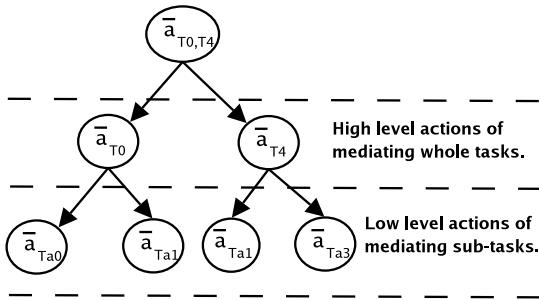
We are interested in evaluating the performance of optimal policies under different termination schemes, with and without the wait operator, and under different distributions of loads. The experimental settings are as follows. The scenario described in Section 2 and pictured in Figure 2 is used. Data are averaged over 10 simulation runs, each consisting of 100,000 episodes (trials). The duration of each subtask is stochastic. With probability  $p_i$ , agent  $i$  will finish

---

**Algorithm 3:** generateActions

---

- 3.1 **begin**
  - 3.2 Every available task (at the current time)  $T_i$  has a set of decompositions  $D_i$ . Each decomposition  $d^j \in D_i$  is a set of subtasks  $\{T_1, \dots, T_n\}$  (In the simple example in Section 2,  $T_0$  has  $D_0 = \{d_0\}$ , i.e., single decomposition, where  $d_0 = \{T_{a_0}, T_{a_1}\}$ )
  - 3.3 A set of *agent-subtask assignments*,  $A_{d_j}$ , is generated for every decomposition. Each subtask assignment  $\bar{a} \in A_{d_j}$  assigns every subtask in the decomposition  $d_j$  to a neighboring agent, i.e.,  $a = \{\langle T_1, a_1 \rangle, \dots, \langle T_n, a_n \rangle\}$ , where  $a_i$  here refers to an agent (in the example in the previous step,  $A_{d_0} = \{\bar{a}_0\}$  because every subtask can be assigned to only one agent, and  $\bar{a}_0 = \{\langle T_{a_0}, a_0 \rangle, \langle T_{a_1}, a_1 \rangle\}$ )
  - 3.4  $A \leftarrow \cup_{d_j} A_{d_j}$
  - 3.5 **end**
- 



**Figure 4:** The hierarchy of the joint action  $\bar{a}_{T_0, T_4}$

its subtask at this time step, and with probability  $1 - p_i$  it will not. Agents  $a_0$  and  $a_2$  have  $p = 0.2$  while  $a_1$  and  $a_3$  have  $p = 0.8$  (i.e. faster agents). Any subtask can take at most 10 time steps. Tasks, arriving to the mediator at time  $t$ , are selected randomly from all combinations of two subtasks (i.e. requiring two agents). Thus we have six different tasks shown in table 1.

When any subtask is achieved the mediator gets a reward of -1 (cost). The *wait* operator has a reward of zero. When *all* subtasks of any task are achieved the mediator gets a reward of 16. Following our framework, contracting a task is modeled as a joint action consisting of two sub-actions. Each sub-action corresponds to assigning a subtask to an agent. Thus, the mediator has at its disposal 6 joint actions corresponding to table 1, as well as the *wait* operator which can be either enabled or disabled (described shortly).

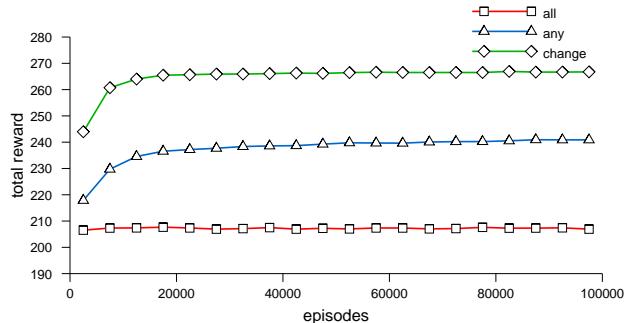
Tasks	$a_0$	$a_1$	$a_2$	$a_3$
$T_0$	*	*		
$T_1$	*		*	
$T_2$	*			*
$T_3$		*	*	
$T_4$		*		*
$T_5$			*	*

**Table 1:** Types of Tasks

Although not required by our model, we deliberately chose that all tasks and subtasks have the same reward so that the only criteria for evaluating a “good” task is its duration. Task  $T_4$  is on average the shortest as it consists of two subtasks that can be achieved by the two fast agents.  $T_4$  is available at any time with probability  $p_{T_4}$ . Any other task is available with probability 0.2.

Defining the state strikes a balance between optimality and speed of convergence. State consists of two parts. The first is the task currently being mediated, which is null if the mediator is idle. The second part is the actual progress of each action (i.e. the mediation process of each subtask). For simplicity our implementation uses the time elapsed while mediating each subtask. Although this state is not the most compact but it allows us to measure the trade-off between the optimality and the speed of convergence of different termination schemes. Note that in Q-learning, an agent learns only about states and actions the agent actually encounters. So for example, using the  $\tau_{all}$  termination scheme, an agent will only learn about the first part of the state. The second part of the state is irrelevant because an agent makes its decision only when all subtasks terminate (either successfully or unsuccessfully). Using  $\tau_{any}$  an agent will witness more states but not all of the states (e.g. at least one action must have terminated). An agent using  $\tau_{change}$  can potentially encounter all states and hence need to learn about all of them.

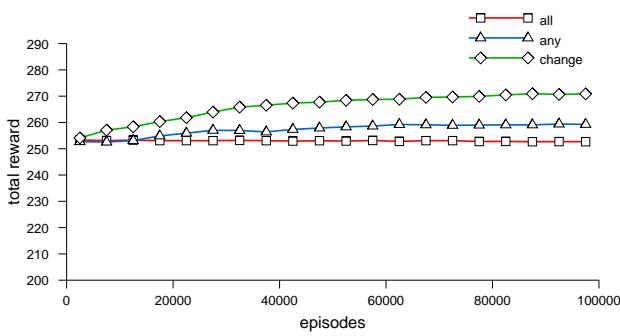
Figure 5 shows the performance of different termination schemes when  $p_{T_4} = 0.6$  and the wait operator is disabled. As shown, the policy learned using  $\tau_{change}$  clearly outperforms those learned using  $\tau_{any}$  or  $\tau_{all}$ . It also converges faster. This was initially surprising as we expected the  $\tau_{change}$  policy to take more time to converge as it tunes the policy with finer granularity. But the reason  $\tau_{change}$  was able to quickly surpass the performance of  $\tau_{any}$  and  $\tau_{all}$  is that the short task  $T_4$  appears so frequently. This gives clear advantage to  $\tau_{change}$  which can interrupt an ongoing task (joint action) even before any subtask (single action) terminates.



**Figure 5:** The performance of different termination schemes when  $p_{T_4} = 0.6$  and the wait operator is disabled

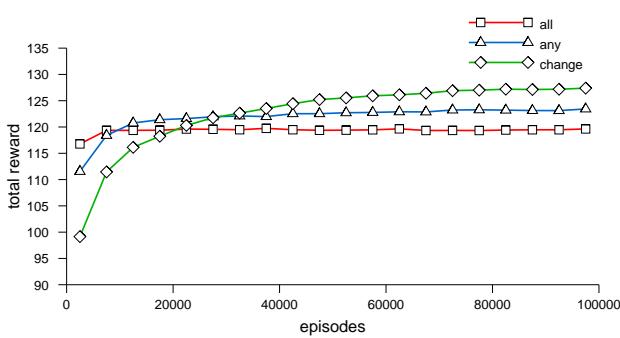
Figure 6 shows the performance of different termination schemes when  $p_{T_4} = 0.6$  and the wait operator is enabled. The wait operator significantly improves the performance of the  $\tau_{all}$  policy and considerably improves the  $\tau_{any}$  policy while slightly improves the performance of the  $\tau_{change}$ . With  $\tau_{all}$  the mediator learns to wait for  $T_4$  type of tasks. The same is true for  $\tau_{any}$  as it can decide to wait if  $T_4$  is not available. However, for  $\tau_{change}$  having the wait operator does not add much (as the mediator can interrupt a task taking too long anyway if there is a better task). In contrast, adding

*wait* as an additional action slows convergence as the agent need to learn about it as well.



**Figure 6:** The performance of different termination schemes when  $p_{T_4} = 0.6$  and the wait operator is enabled

Figure 7 compares the performance when  $p_{T_4} = 0.1$  and the wait operator enabled. Though the policy using  $\tau_{change}$  still outperforms other termination schemes, its convergence is considerably slower than the other two. The reason is that there are more decision epochs and states visited by  $\tau_{change}$  than the other two termination schemes (which is the same reason why  $\tau_{any}$  converges slower than  $\tau_{all}$ ), while the benefit of  $\tau_{change}$  in interrupting an ongoing task is reduced due to lower  $p_{T_4}$ .

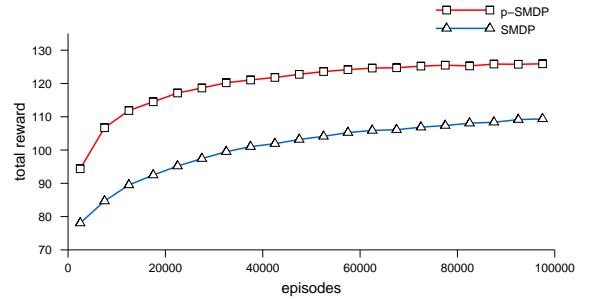


**Figure 7:** The performance of different termination schemes when  $p_{T_4} = 0.1$  and the wait operator is enabled

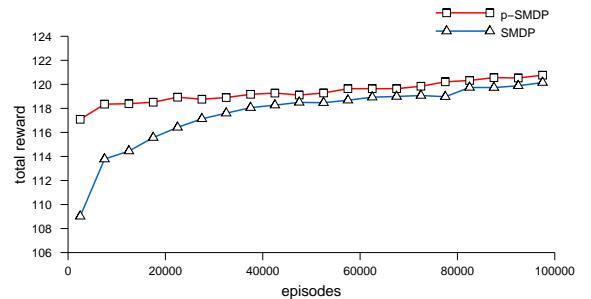
Figures 8 and 9 verify the theoretical results in Section 4. The convergence of the learned policies with the set of available actions factored out of the state (i.e. relying on  $\wp$ -SMDP model) is considerably faster than the learned policies with the set of available actions being part of the state. The reason is the size of the state space. When the set of available actions is part of the state, the mediator needs to learn about 32 times the number of states when the set of available actions is factored out. This slows convergence considerably, especially in the case of  $\tau_{change}$  (Figure 8) because the size of the state space, even when the of set available actions are factored out, is high.

## 7. RELATED WORK

The authors in [4] used an MDP model for the decision of a meta-controller. This meta-controller chose among high-level actions of



**Figure 8:** The performance of the policy learned using  $\tau_{change}$  with available actions as part of the state (SMDP) and factored out of the state ( $\wp$ -SMDP).



**Figure 9:** The performance of the policy learned using  $\tau_{all}$  with available actions as part of the state (SMDP) and factored out of the state ( $\wp$ -SMDP).

the agent including: choosing a task to pursue among a set of available tasks, the detail with which to schedule a task, when to abort a task, when to start and stop a negotiation (mediation) with another agent. Though their decision problem is similar to ours, the approach taken is considerably different. They used the traditional MDP model, which is not scalable in the number of available tasks. Therefore, they had to limit the number of available actions at any time. Our  $\wp$ -SMDP model scales well with the number of available tasks. Also their model only handled serialized actions, while our model deals with concurrent actions.

In [2], a mediator serially allocate tasks to agents. They used an MDP model where actions are agent-task assignments. This again differs from our concurrent allocation model, with different action durations. They also assumed that the set of tasks to be allocated are *fixed* and arrive in fixed order, while we assume tasks arrive stochastically. They also assumed agents with homogeneous capabilities, while our model supports heterogeneous agents.

The work in [1] modeled the resource allocation problem as a constrained MDP, or CMDP. A CMDP is an MDP augmented with a set of (resource) constraints. They did not support infinite horizon decision processes, unlike our model. The set of actions were assumed fixed and the policy was serial. They also used an offline algorithm which solved the problem assuming the transition probabilities are known.

## 8. CONCLUSION AND FUTURE WORK

We present in this work a single coherent Semi-MDP model of the challenging mediator decision problem. We propose the  $\wp$ -SMDP model, which is a generalization of the SMDP model, and we prove that  $\wp$ -SMDP model can lead to *exponential* savings in the memory and time required to learn an optimal policy. We also extended the concurrent action model with a new termination scheme that results in an optimal policy outperforming all previous termination schemes.

This work lays the basic formal foundation for more interesting questions to be answered in the future. This work's focus is on modeling the decision process of a single mediator. Evaluating the performance of our model and our algorithm in systems with a network of mediators is interesting and an ongoing work. The issues we intend to address include: how the network connections between mediators and agents affect performance, how the load of tasks arriving at one mediator affect the performance at another mediator, and how load balancing can be incorporated in our model.

While our model can represent concurrent mediations as joint actions, the model is exponential in the size of the joint actions. This is clearly not scalable. We are currently investigating under what assumptions our model can be *factored* over tasks and how it affects optimality.

## 9. REFERENCES

- [1] D. Dolgov and E. Durfee. Optimal resource allocation and policy formulation in loosely-coupled markov decision processes. In *In Proceedings of the Fourteenth International Conference on Automated Planning and Scheduling*, 2004.
- [2] H. Hannah and A.-I. Mouaddib. Task selection problem under uncertainty as decision-making. In *Proceedings of the first international joint conference on Autonomous agents and multiagent systems*, 2002.
- [3] M. L. Littman and G. Szepesvari. A generalized reinforcement-learning model: Convergence and applications. In L. Saitta, editor, *Proceedings of the 13th International Conference on Machine Learning (ICML-96)*, pages 310–318, Bari, Italy, 1996. Morgan Kaufmann.
- [4] A. Raja and V. Lesser. Meta-level Reasoning in Deliberative Agents. *Proceedings of the International Conference on Intelligent Agent Technology (IAT 2004)*, September 2004.
- [5] K. Rohanianesh and S. Mahadevan. Learning to take concurrent actions. *Sixteenth International Conference on Neural Information Processing Systems (NIPS)*, 2003.
- [6] R. G. Smith. The contract net protocol: High level communication and control in a distributed problem solver. *IEEE Transactions on Computers*, 29(12):1104–1113, 1980.
- [7] R. Sutton and A. Barto. *Reinforcement Learning: An Introduction*. MIT Press, 1999.
- [8] R. S. Sutton, D. Precup, and S. P. Singh. Between MDPs and semi-MDPs: A framework for temporal abstraction in reinforcement learning. *Artificial Intelligence*, 112(1-2):181–211, 1999.